

3.16.1. Derived Rule Problems

For each of the following deductive systems, show that one or more of our deductive rules are **derived rules** in that system (by constructing a deduction of that rule, in that system).

1. Show that the rule **MP** can be treated as a derived rule in the Chapter Three deductive system, by providing a deduction of the argument “ $(P \rightarrow Q) \cdot P \therefore Q$ ” that doesn’t use MP.
2. Show that the rule **MT** can be treated as a derived rule in the Chapter Three deductive system, by providing a deduction of the argument “ $(P \rightarrow Q) \cdot \sim Q \therefore \sim P$ ” that doesn’t use MT.
3. Show that the rule **R** can be treated as a derived rule in the Chapter Three deductive system, by providing a deduction of the argument “ $P \therefore P$ ” **using only CD and MP**. (*From Kalish, Montague, and Mar 1980: 48 #38*)
4. Show that the rule $\sim -$ can be treated as a derived rule in the Chapter Three deductive system, by providing a deduction of the argument “ $\sim \sim P \therefore P$ ” **using only ID, CD, and MP**. (*From Kalish, Montague, and Mar 1980: 48 #38*)

5. The deductive system **DS1** has only ID, plus the rules $\wedge-$, $\wedge+$, $\vee-$, $\vee+$, and “Negated Conditional” (“ $\sim\rightarrow$ ”).

Negated Conditional ($\sim\rightarrow$)

$$\frac{\sim(\bullet \rightarrow \blacktriangle)}{\therefore (\bullet \wedge \sim\blacktriangle)} \qquad \frac{(\bullet \wedge \sim\blacktriangle)}{\therefore \sim(\bullet \rightarrow \blacktriangle)}$$

Thanks to $\sim\rightarrow$, any **CD** from our system of deduction can be converted into an ID in DS1.¹

5a. Show that **MT** is a derived rule in DS1.

5b. Show that **MP** is a derived rule in DS1.

6. The deductive system **DS2** is like our Chapter Three deductive system except that it lacks the rule $\vee-$ and instead has the rule **Separation of Cases (SC)**.

Separation of Cases (SC)

$$\frac{\begin{array}{c} (\bullet \vee \heartsuit) \\ (\bullet \rightarrow \blacktriangle) \\ (\heartsuit \rightarrow \blacktriangle) \end{array}}{\blacktriangle}$$

¹ Because DS1 has **ID** and the rules $\vee-$ and $\vee+$, the rules R , $\sim-$, and $\sim+$ are already derived rules in this system – as shown in 2.40.1 Problems 2, 3, and 5.

So the following argument is an example of (SC).

$$\begin{array}{l}
 1. (P \vee Q) \\
 2. (P \rightarrow R) \\
 3. (Q \rightarrow R) \\
 \hline
 \therefore R
 \end{array}$$

6a. Show that the above argument is deducible in the Chapter Three deductive system (and hence that **SC** can act as a derived rule in the Chapter Three system).

6b. Show that the following argument is deducible using only **CD**, **ID**, **R**, and **SC** (and hence that \vee – is a derived rule in **DS2**).

$$\begin{array}{l}
 1. (P \vee Q) \\
 2. \sim P \\
 \hline
 \therefore Q
 \end{array}$$

6c. Show that the following argument is deducible using only **CD**, **ID**, \wedge –, **R**, and **SC**.

$$\begin{array}{l}
 1. (\sim P \wedge \sim Q) \\
 \hline
 \therefore \sim(P \vee Q)
 \end{array}$$

7. The following argument is an instance of the rule **Double Disjunction** (DD), discussed in Chapter Two.²

$$\begin{array}{l} 1. (P \vee Q) \\ 2. (P \vee \sim Q) \\ \hline \therefore P \end{array}$$

Provide a deduction of the above argument using only $\sim+$, $\vee-$, **SC**, and **CD**.

8. Note that the argument “ $(P \rightarrow (P \rightarrow Q)) \therefore (P \rightarrow Q)$ ” is valid. We can use the rule Separation of Cases to **explain** the validity of this rule **deductively**: from the sentence “ $(P \rightarrow (P \rightarrow Q))$,” when accompanied by the sentences “ $(P \vee \sim P)$ ” and “ $(\sim P \rightarrow (P \rightarrow Q))$,” the conclusion “ $(P \rightarrow Q)$ ” follows as an instance of Separation of Cases.

$$\begin{array}{l} 1. (P \vee \sim P) \\ 2. (P \rightarrow (P \rightarrow Q)) \\ 3. (\sim P \rightarrow (P \rightarrow Q)) \\ \hline \therefore (P \rightarrow Q) \quad 1, 2, 3, \text{SC} \end{array}$$

But the other two premises, “ $(P \vee \sim P)$ ” and “ $(\sim P \rightarrow (P \rightarrow Q))$,” are **theorems** provable without premises³ – in effect, built in the deductive system. In that sense, from the deductive system alone “ $(P \rightarrow (P \rightarrow Q))$ ” entails “ $(P \rightarrow Q)$ ” by Separation of Cases.

For each of the following arguments, use **Separation of Cases** plus theorems to deductively explain the validity of that argument.

8a. $(P \vee Q) \cdot (P \rightarrow Q) \therefore Q$

8b. $(\sim P \rightarrow P) \therefore P$

8c. $(P \rightarrow \sim(P \wedge Q)) \therefore \sim(P \wedge Q)$

² In 2.40.1 Problem 8

³ “ $(P \vee \sim P)$ ” is theorem **T2.2** from 2.41.1, and “ $(\sim P \rightarrow (P \rightarrow Q))$ ” is theorem **T3.3b** from 3.13.1 C.

8d. $(P \rightarrow (Q \rightarrow \sim P)) \therefore (Q \rightarrow \sim P)$

8e. $((P \rightarrow Q) \rightarrow P) \therefore P$

9. We can derive a number of rules using the rule **Contradictory Consequents (CC)**.

Contradictory Consequents (CC)

$$\frac{\begin{array}{l} (\bullet \rightarrow \blacktriangle) \\ (\bullet \rightarrow \sim \blacktriangle) \end{array}}{\therefore \sim \bullet}$$

For instance, MT becomes a derived rule (using only CD, R, and CC).

- | | | |
|----|--------------------------|-------------------------------|
| 1. | $(P \rightarrow Q)$ | |
| 2. | $\sim Q$ | |
| | | Get: $\sim P$ |
| | | Get: $(P \rightarrow \sim Q)$ |
| 3. | P | ACD |
| 4. | $\sim Q$ | 2, R |
| 5. | $(P \rightarrow \sim Q)$ | 3, 4, CD |
| 6. | $\sim P$ | 1, 5, CC |

Provide **deductions** of each of the following arguments (using only the rules and forms of deduction specified).

9a. $P \therefore \sim \sim P$ (using only CD, R, and CC)

9b. $(P \rightarrow Q) \cdot P \therefore Q$ (using only ID, CD, R, and CC)

10. We can derive a number of rules using a variant of the **Contradictory Consequents rule, (CC*)**.

Contradictory Consequents* (CC*)⁴

$$\frac{\begin{array}{l} (\sim \bullet \rightarrow \blacktriangle) \\ (\sim \bullet \rightarrow \sim \blacktriangle) \end{array}}{\therefore \bullet}$$

Provide **deductions** of each of the following arguments (using only the rules and forms of deduction specified).

10a. $\sim\sim P \therefore P$ (using only CD, R, and CC*)

10b. $(P \rightarrow Q) \cdot P \therefore Q$ (using only CD, R, MT, and CC*)

10c. $(P \rightarrow Q) \cdot P \therefore Q$ (using only CD, R, CC, and CC*)

10d. $(\sim P \rightarrow \sim Q) \cdot P \therefore Q$ (using only CD, R, and CC*)

⁴ This is the conditional analogue of the Chapter Two derived rule **Double Disjunction** (discussed in 2.40.1 Problem 8).

11. We can derive a number of rules using a **variant of Modus Tollens**, (MT*).

Modus Tollens* (MT*)

$$\frac{(\sim \bullet \rightarrow \sim \blacktriangle) \quad \blacktriangle}{\therefore \bullet}$$

Provide **deductions** of each of the following arguments (using only the rules and forms of deduction specified). (*11a-11c adapted from Kalish, Montague, and Mar 1980: 44 #39*)

11a. $P \therefore P$ (using only CD and MT*)

11b. $\sim\sim P \therefore P$ (using only CD and MT*)

11c. $P \therefore \sim\sim P$ (using only CD and MT*)

11d. $(\sim P \rightarrow \sim Q) \cdot P \therefore Q$ (using only MP, $\sim+$, and the following rule of **Contraposition**)

Contraposition (CP)

$$\frac{(\sim \bullet \rightarrow \sim \blacktriangle)}{\therefore (\blacktriangle \rightarrow \bullet)}$$

11e. $(\sim P \rightarrow \sim Q) \therefore (Q \rightarrow P)$ (using only CD, R, and CC*)

11f. $(Q \rightarrow P) \therefore (\sim P \rightarrow \sim Q)$ (using only CD, R, and CC)